**Instructions:** In-class exercises are meant to introduce you to a new topic and provide some practice with the new topic. Work in a team of up to 4 people to complete this exercise. You can work simultaneously on the problems, or work separate and then check your answers with each other. Turn in <u>one</u> copy of the exercise per group.

Names:

# Boolean Algebra

## Logic, Sets, and Boolean Algebra

Logic  $(p \land q)$ , sets  $(P \cap Q)$ , and now boolean algebra  $(p \cdot q)$  have a lot in common with each other. In fact, sometimes it can be useful to convert a problem from one type to another in order to learn more about a problem. Here are the "translations":

|                 | Logic             | Sets                  | Boolean<br>Algebra |
|-----------------|-------------------|-----------------------|--------------------|
| Variables       | p, q, r           | A, B, C               | a, b, c            |
| "and" operation | $\wedge$          | $\cap$                |                    |
| "or" operation  | V                 | U                     | +                  |
| "not" operation | _                 | ,                     | 1                  |
| "-" operation   | $a \wedge \neg b$ | A - B                 | $a \cdot b'$       |
| Special         | Tautology         | Universal set U       | 1                  |
|                 | Contradiction     | Empty set $\emptyset$ | 0                  |

#### Example:

Rephrase the following Logic operation using Set and Boolean Algebra notations:  $(p \wedge q) \vee r$ 

- Logic:  $(p \land q) \lor r$
- Sets:  $(P \cap Q) \cup R$
- Boolean algebra:  $(p \cdot q) + r$

#### Question 1

Rewrite the following using Boolean Algebra notation:

a.  $p \land q$ b.  $p \lor q$ c.  $\neg p$ d.  $(p \land \neg q) \lor p$ e.  $\neg (\neg p)$ f.  $(p \land \neg q) \lor p \equiv p$ g. (A - B)h.  $A' \cup (A \cap B)$ i.  $(A - B)' = A' \cup (A \cap B)$ 

# Boolean Algebra properties

| Boolean Algebra Properties <sup>a</sup>                    |  |  |  |  |
|--|--|--|--|--|
| Commutative  | $a \cdot b = b \cdot a$  | a+b=b+a  |  |  |
| Associative  | $(a \cdot b) \cdot c = a \cdot (b \cdot c)$                                  | (a+b) + c = a + (b+c)                          |  |  |
| Distributive   | $\begin{aligned} a \cdot (b+c) \\ = (a \cdot b) + (a \cdot c) \end{aligned}$ | $a + (b \cdot c)$<br>= $(a + b) \cdot (a + c)$ |  |  |
| Identity   | $a \cdot 1 = a$  | a + 0 = a                                      |  |  |
| Negation   | a + a' = 1   | $a \cdot a' = 0$                               |  |  |
| Double negative  | (a')' = a  |  |  |  |
| Idempotent   | $a \cdot a = a$  | a + a = a                                      |  |  |
| DeMorgan's laws  | $(a \cdot b)' = a' + b'$   | $(a+b)' = a' \cdot b'$                         |  |  |
| Universal bound  | a + 1 = 1  | $a \cdot 0 = 0$                                |  |  |
| Absorption   | $a \cdot (a+b) = a$  | $a + (a \cdot b) = a$                          |  |  |
| Complements<br>of 1 and 0                                  | 1' = 0   | 0' = 1   |  |  |
| <sup>a</sup> From Discrete Mathematics, Ensley and Crawley |  |  |  |  |

Examples <sup>a</sup> **Example 1:** Simplify  $(a + 1) \cdot (a + 0)$  $(a+1)\cdot(a+0)$  $= 1 \cdot (a+0)$ universal bound  $= (a+0) \cdot 1$ commutative = a + 0identity = aidentity **Example 2:** Simplify  $a \cdot (a' + b)$  $a \cdot (a'+b)$  $= a \cdot a' + a \cdot b$ distributive  $= 0 + a \cdot b$ negation  $= a \cdot b + 0$ commutative  $= a \cdot b$ identity Show that  $a \cdot b + b \cdot c = (a + c) \cdot b$ Example 3:  $a \cdot b + b \cdot c$  $= b \cdot a + b \cdot c$  commutative  $= b \cdot (a+c)$ distributive  $= (a+c) \cdot b$ commutative **Example 4:** Show that if a' + b = 1 then  $a \cdot b' = 0$ . a' + b $= (a')' \cdot b'$ double negative = (a' + b)'DeMorgan =(1)'Since a' + b = 1= 0Complements <sup>a</sup>From Discrete Mathematics, Ensley & Crawley, 3.4

### Question 2

Simplify the following equations using the Boolean Algebra properties.

a. ab + ab'

b.  $cd \cdot c'd'$ 

c.  $e \cdot (e+f)$ 

d.  $g \cdot (h+i)$ 

e.  $j + (k \cdot l)$ 

f. m + (no + no')

#### Example

Using the Boolean Algebra properties, transform the left-hand side of<br/>each equation to the right-hand side.<br/>Show that ba + ba' = bStep<br/> $ba + ba' \rightarrow b \cdot (a + a')$ 1. Distributive  $a \cdot (b + c) = (ab) + (ac)$ Step<br/> $ba + ba' \rightarrow b \cdot (a + a')$ 2. Negation a + a' = 1 $b \cdot (a + a') \rightarrow b \cdot (1)$ 3. Identity  $a \cdot 1 = a$  $b \cdot (1) \rightarrow b \cdot 1 = b$ 

Therefore ba + ba' = b

#### Question 3

Using the Boolean Algebra properties, transform the left-hand side of each equation to the right-hand side.

c. x'yz + x'y'z + xyz' + xy'z' = xz' + x'z(Group up two terms at a time: xyz' + xy'z' and x'yz + x'y'z.)

d. xyz + xyz' + x'yz + x'yz' = y

## Logic Circuits



## Question 4



Write out the Boolean expression that describes each diagram:

### Question 5

Draw a circuit diagram for the following Boolean expressions: a. a + b' | b.  $a' \cdot b'$  | c.  $a + (b \cdot c)$